

## Free extensions

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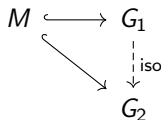
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# Abelian groups and commutative monoids

## Fact

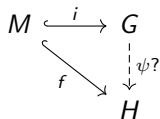
For each commutative monoid  $M$ , there is (up to isomorphism) at most one embedding  $M \hookrightarrow G$  of  $M$  into an Abelian group  $G$  whose image generates  $G$  as a group.

Uniqueness up to isomorphism means:



Example:  $\mathbb{N} \hookrightarrow \mathbb{Z}$  (wrt addition).

Sketch of proof. Let  $i$  and  $f$  be injective monoid homomorphisms with generating images. Wlog,  $i: M \subseteq G$ .



Every  $z \in G$  is a difference  $z = x - y$  of elements of  $M$ .

Set  $\psi(z) := f(x) - f(y)$ .

It is well-defined:

$$\begin{aligned}
 x - y = x' - y' &\iff x + y' = x' + y \\
 &\implies f(x + y') = f(x' + y) \\
 &\iff f(x) + f(y') = f(x') + f(y) \\
 &\iff f(x) - f(y) = f(x') - f(y').
 \end{aligned}$$

Further, one proves that  $\psi$  is a group isomorphism that extends  $f$ .

## Key fact used

For every Abelian group  $G$  and all  $x, y, x', y' \in G$ ,

$$x - y = x' - y' \iff x + y' = x' + y.$$

## Fact

Any equation between two terms in the language of Abelian groups is equivalent to an equation in the language of monoids.

Example:  $-x + y - z = x \iff y = x + x + z.$

# Groups and monoids

For arbitrary (i.e. non-necessarily commutative) monoids/groups, we do **not** have this uniqueness of embeddings.

E.g.: the monoid  $\{x, y, z\}^*$  of words on three letters has two non-isomorphic generating embeddings into groups:

► into the free group

$\text{Free}(\{x, y, z\})$

$\{x, y, z\}^* \longrightarrow \text{Free}(\{z, y, z\})$

$x \longmapsto x$

$y \longmapsto y$

$z \longmapsto z.$

► into the free group

$\text{Free}(\{x, y\})$

$\{x, y, z\}^* \longrightarrow \text{Free}(\{x, y\})$

$x \longmapsto x$

$y \longmapsto y$

$z \longmapsto xy^{-1}x.$

For groups, the equation  $z = xy^{-1}x$  cannot be expressed via an equation in the language of monoids.

## Main result

Our main result is an equivalence between two conditions: the “unique embeddability property” and the “expressibility property”.



Recall that a *prevariety* (a.k.a. *SP-class*) is a class of similar algebras closed under subalgebras and products.

Examples: any variety, any quasi-variety.

## Setting

- ▶ An algebraic language  $\mathcal{L}_+$  and a sublanguage  $\mathcal{L}_- \subseteq \mathcal{L}_+$ .
- ▶ Two prevarieties  $\mathcal{V}_+$  and  $\mathcal{V}_-$  for  $\mathcal{L}_+$  and  $\mathcal{L}_-$  respectively, such that all  $\mathcal{L}_-$ -reducts of algebras in  $\mathcal{V}_+$  belong to  $\mathcal{V}_-$ .

For example:

1.  $\mathcal{V}_+ = \{\text{Abelian groups}\}$ ,  $\mathcal{V}_- = \{\text{commutative monoids}\}$ .
2.  $\mathcal{V}_+ = \{\text{groups}\}$ ,  $\mathcal{V}_- = \{\text{monoids}\}$ .

Recall: Every commutative monoid has at most one embedding into an Abelian group with generating image.

## Definition

*Unique embeddability property* := for all  $A \in \mathcal{V}_-$ ,  $B, C \in \mathcal{V}_+$  and injective  $\mathcal{V}_-$ -homomorphisms  $f: A \hookrightarrow B$  and  $g: A \hookrightarrow C$  whose images  $\mathcal{V}_+$ -generate  $B$  and  $C$  respectively, there is a  $\mathcal{V}_+$ -isomorphism  $h: B \rightarrow C$  making the following diagram commute.

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ & \searrow g & \downarrow \exists! h \\ & & C \end{array}$$

For  $\mathcal{V}_+ = \{\text{Abelian groups}\}$  and  $\mathcal{V}_- = \{\text{commutative monoids}\}$ : the unique embeddability property holds.

For  $\mathcal{V}_+ = \{\text{groups}\}$  and  $\mathcal{V}_- = \{\text{monoids}\}$ : the unique embeddability property does not hold.

An equivalent formulation is: every injective  $\mathcal{V}_-$ -homomorphism  $A \hookrightarrow B$  ( $A \in \mathcal{V}_-, B \in \mathcal{V}_+$ ) whose image  $\mathcal{V}_+$ -generate  $B$  is a *free extension*, i.e. the free way in which  $A$  embeds into an algebra in  $\mathcal{V}_+$ . In other words, it is (up to an iso) the component at  $A$  of the unit of the adjunction that has the forgetful functor  $\mathcal{V}_+ \rightarrow \mathcal{V}_-$  as a right adjoint.

In short: “injective + generating  $\Rightarrow$  free”.

## Definition

*Expressibility property* := every equation in  $\mathcal{L}_+$  is equivalent to a system of equations in  $\mathcal{L}_-$ .

I.e.: for each pair  $(\sigma(x_1, \dots, x_n), \rho(x_1, \dots, x_n))$  of terms in  $\mathcal{L}_+$ , there is a (finite) set of pairs  $(\alpha_i(x_1, \dots, x_n), \beta_i(x_1, \dots, x_n))_i$  of terms in  $\mathcal{L}_-$  s.t., for all  $A \in \mathcal{V}_+$  and  $x_1, \dots, x_n \in A$ ,

$$\sigma(x_1, \dots, x_n) = \rho(x_1, \dots, x_n) \Leftrightarrow \forall i \alpha_i(x_1, \dots, x_n) = \beta_i(x_1, \dots, x_n).$$

For  $\mathcal{V}_+ = \{\text{Abelian groups}\}$  and  $\mathcal{V}_- = \{\text{commutative monoids}\}$ : it holds.  
E.g.  $x - y = x' - y'$  iff  $x + y' = x' + y$ .

For  $\mathcal{V}_+ = \{\text{groups}\}$  and  $\mathcal{V}_- = \{\text{monoids}\}$ : it does not hold. E.g.:  
 $z = xy^{-1}x$ .

## Main theorem

Unique embeddability property  $\iff$  expressibility property.

## Examples

# Rational vector spaces and Abelian groups

$\mathcal{V}_+ := \{\text{rational vector spaces}\}$ ,  $\mathcal{V}_- := \{\text{Abelian groups}\}$ .

True or false? (Unique embeddability property)

For each Abelian group  $G$ , there is (up to isomorphism) at most one embedding (of Abelian groups)  $G \hookrightarrow V$  into a rational vector space whose image generates  $V$  (as a rational vector space).

(Example:  $\mathbb{Z} \hookrightarrow \mathbb{Q}$ .)

Expressibility property?  $\frac{1}{2}x + y = \frac{2}{3}z$  is equivalent to  $3x + 6y = 4z$ .

Fact (Expressibility property)

Every equation in the language of  $\mathbb{Q}$ -vector spaces is equivalent to an equation in the language of Abelian groups.

Answer: true.

# Complex and real vector spaces

$\mathcal{V}_+ = \{\text{complex vector spaces}\}$ ,  $\mathcal{V}_- = \{\text{real vector spaces}\}$ .

## True or false? (Unique embeddability property)

For each real vector space  $V$ , there is (up to isomorphism) at most one embedding (of Abelian groups)  $V \hookrightarrow W$  into a complex vector space whose image generates  $W$  (as a complex vector space).

Expressibility property? The equation  $x = iy$  does **not** seem to be equivalent to a system of equations in the language of real vector spaces.

The embeddings  $f: \mathbb{R} \times \mathbb{R} \hookrightarrow \mathbb{C} \times \mathbb{C}$  and  $g: \mathbb{R} \times \mathbb{R} \hookrightarrow \mathbb{C}$  are non-isomorphic.  $f(0, 1) \neq if(1, 0)$ ,  $g(0, 1) = ig(1, 0)$ .

Answer: false.

In turn, we can then infer that the equation  $x = iy$  is not equivalent to an equation in the language of real vector space.



# Abelian $\ell$ -groups and monoids

$\mathcal{V}_+$  = Abelian lattice-ordered groups.

$\mathcal{V}_-$  = commutative lattice-ordered monoids.

The expressibility and the unique embeddability properties hold.

Example:  $x \vee (z - y) = x - z$  is equivalent to  
 $(x + y + z) \vee (z + z) = x + y$ .

# Boolean algebras and lattices

$\mathcal{V}_+ = \{\text{Boolean algebras}\}$ ,  $\mathcal{V}_- = \{\text{bounded distributive lattices}\}$ .

## Expressibility property

Every equation in the language of Boolean algebras is equivalent to a system of equations in the language of bounded distributive lattices.

Example:  $x \vee \neg y = \neg z$  is equivalent to 
$$\begin{cases} x \wedge z = 0; \\ x \vee y = x \vee z. \end{cases}$$

## Unique embeddability property

For each bounded distributive lattice  $L$ , there is (up to isomorphism) at most one embedding  $L \hookrightarrow B$  into a Boolean algebra whose image generates  $B$  (as a Boolean algebra).

This embedding is called the *free Boolean extension* of  $L$ .

# MV-algebras and positive MV-algebras

$\mathcal{V}_+$  = MV-algebras.

$\mathcal{V}_-$  = positive MV-algebras (i.e.  $\{\oplus, \odot, \vee, \wedge, 0, 1\}$ -subreducts of MV-algebras).

Both the expressibility and the unique embeddability properties hold.

# Modal algebras and positive modal algebras

$\mathcal{V}_+$  = Modal algebras.

$\mathcal{V}_-$  = Positive modal algebras.

The unique embeddability and the expressibility properties do *not* hold.

Example: the equation  $\Diamond(x \wedge \neg y) = 0$  is not expressible via a system of equations in the language of positive modal algebras.

## Summing it up and future directions

# Summing it up and future directions

## Main theorem

Unique embeddability property  $\iff$  expressibility property.

This was established in the setting of prevarieties.

Future direction (from a suggestion of a referee): does the equivalence extend beyond algebraic structures to a general model-theoretic setting?

Thank you!