

Natural dualities

Marco Abbadini. University of Birmingham, UK.

Joint work (in progress) with Adam Přenosil.

Southern and Midlands Logic Seminar, 13 December 2023.

1936: Stone's representation theorem for Boolean algebras:

Every Boolean algebra is isomorphic to a field of subsets of some set.

It connects syntax and semantics: each formula of classical propositional logic is interpreted as a subset of the set of possible worlds, where

- ▶ logical “or” \leftrightarrow union of sets of worlds,
- ▶ logical “and” \leftrightarrow intersection,
- ▶ logical “negation” \leftrightarrow complementation.

Given a Boolean algebra \mathbf{B} , Stone equips the set $X_{\mathbf{B}}$ of homomorphisms $\mathbf{B} \rightarrow \mathbf{2}$ with an appropriate topology and shows

$$\mathbf{B} \cong \{\text{clopen (= closed + open) subsets of } X_{\mathbf{B}}\}.$$

One-to-one correspondence between Boolean algebras and Stone spaces (a.k.a. profinite spaces or Boolean spaces).

Stone space := compact Hausdorff spaces where distinct points are separated by a clopen (= closed and open) set.

E.g.: finite discrete spaces, one-point compactifications of discrete spaces.

Stone duality

In fact, Stone proved much more than a representation theorem: there is a duality (= dual categorical equivalence) between the category of Boolean algebras and the category of Stone spaces.

$$\text{Bool} \cong \text{Stone}^{\text{op}}.$$

More information (quotient of algebras $A \twoheadrightarrow B$)

=

Fewer possible worlds (inclusion of spaces $X_A \hookrightarrow X_B$).

Less propositions (inclusion of algebras $B \hookrightarrow A$)

=

Collapse of possible worlds (quotient of spaces $X_B \leftarrow X_A$).

Some good things of Stone duality

Algebraic questions in Bool can be answered by translating them into (often simpler) questions in Stone.

1. While a coproduct $\mathbf{A} + \mathbf{B}$ of Boolean algebras is often difficult to describe, its dual, $X_{\mathbf{A}+\mathbf{B}}$, is simply the cartesian product $X_{\mathbf{A}} \times X_{\mathbf{B}}$.
2. The free Boolean algebra (= Lindenbaum-Tarski algebras) on κ generators is easily described via its dual: 2^κ .
3. Congruences correspond to closed subspaces.
Example: a Boolean algebra of 2^n elements has exactly 2^n congruences.

Some good things of Stone duality

The duality is logarithmic: a product $\mathbf{A} \times \mathbf{B}$ of Boolean algebras corresponds to a sum of the corresponding spaces: $X_{\mathbf{A} \times \mathbf{B}} \cong X_{\mathbf{A}} + X_{\mathbf{B}}$.

For example, the dual of a Boolean algebra of $16 = 2 \times 2 \times 2 \times 2$ elements is the discrete space of $4 = 1 + 1 + 1 + 1$ elements.

Priestley duality

[1969, Priestley]: the category of bounded distributive lattices is dually equivalent to the category of Priestley spaces.

$$\text{BDL} \cong \text{Priestley}^{\text{op}}.$$

Priestley space := Stone space equipped with a partial order such that if $x \not\leq y$ then there is a clopen upset U such that $x \in U$ and $y \notin U$.

Every bounded distributive lattice is isomorphic to the lattice of clopen upsets of a Priestley space.

New classes of algebras

To study some nonclassical logics, new classes of algebras (other than Boolean algebras) were introduced.

E.g.: Kleene (1938) introduced a three-valued logic which replaces $\{0, 1\}$ with $\{0, U, 1\}$. $\{0, U, 1\}$ is an algebra in the signature $\{0, 1, \vee, \wedge, \neg\}$.

Kleene algebras := algebras of Kleene logic = algebras in the signature $\{0, 1, \vee, \wedge, \neg\}$ that satisfy each equation satisfied by $\{0, U, 1\}$:

- ▶ Bounded distributive lattices.
- ▶ (Double negation law:) $\neg\neg x = x$.
- ▶ (De Morgan laws:) $\neg(x \wedge y) = \neg x \vee \neg y$.
- ▶ $x \wedge \neg x \leq y \vee \neg y$.

In general, we do *not* have $x \wedge \neg x = 0$.

Dualities for many of these classes of algebras were found.

Benefits similar to those for Stone duality apply:

- ▶ uniform representation,
- ▶ coproducts of algebras into products of spaces,
- ▶ free algebras have an easy dual description,
- ▶ congruences correspond to closed substructures,
- ▶ logarithmic property: $X_{\mathbf{A} \times \mathbf{B}} \cong X_{\mathbf{A}} + X_{\mathbf{B}}$.

These dualities follow a similar pattern: there is a special algebra

L

and the class of algebras of interest is

$\text{ISP}(\mathbf{L}) := \{\text{algebras isomorphic to a subalgebra of a power of } \mathbf{L}\}.$

i.e. the class of algebras **A** (in the same language of **L**) whose elements are separated by homomorphisms from **A** to **L**.

Examples:

- ▶ Boolean algebras = $\text{ISP}(\{0, 1\})$ (by the Ultrafilter lemma).
- ▶ Bounded distributive lattices = $\text{ISP}(\{0, 1\})$.
- ▶ Kleene algebras = $\text{ISP}(\{0, U, 1\})$.

Logical reading: **L** is the set of truth values.

Homming to **L** then gives the functor for the duality.

Natural dualities $:=$ dualities induced by homming into a schizophrenic object.

[Davey, 1978], [Davey, Werner, 1983], [Clark, Krauss, 1984]: formulate a general theory to obtain natural dualities for $\mathbb{ISP}(\mathbf{L})$.

Davey & Werner (1983): duality for $\mathbb{ISP}(\mathbf{L})$ where \mathbf{L} is a finite algebra with a near unanimity term. (E.g.: \mathbf{L} has a lattice reduct, i.e. \mathbf{L} has two terms \vee and \wedge that satisfy the axioms of lattices.)

Example: for $\mathbf{L} = \{0, 1\}$ we get Stone/Priestley dualities.

We present the restriction of this theory under some stronger hypotheses on \mathbf{L} :

0. \mathbf{L} is finite.
1. \mathbf{L} has a lattice reduct.
2. \mathbf{L} has two distinct constants and every subalgebra of \mathbf{L} is simple (i.e. the only congruences of a subalgebra \mathbf{A} are the diagonal and $\mathbf{A} \times \mathbf{A}$).
3. For each subalgebra \mathbf{A} of \mathbf{L} , the inclusion $\mathbf{A} \hookrightarrow \mathbf{L}$ is the unique homomorphism from \mathbf{A} to \mathbf{L} .

An algebra in $\mathbf{ISP}(\mathbf{L})$ is represented via a

1. Stone space X , together with
2. a family of “at-most-binary constraints”, i.e., for each subset $I \subseteq X$ of cardinality ≤ 2 , a subalgebra \mathbf{A}_I of \mathbf{L}^I .

satisfying certain compatibility conditions.

Space X + constraints $(\mathbf{A}_I)_I \rightsquigarrow$ Algebra of continuous functions $f: X \rightarrow \mathbf{L}$ (with \mathbf{L} discrete) satisfying all constraints, i.e. such that, for all $I \subseteq X$ with $|I| \leq 2$, $f|_I \in \mathbf{A}_I$.

1. Stone duality for Boolean algebras (1936): $\mathbf{L} = \{0, 1\}$.

Space: a Stone space.

Constraints: “none”, i.e. $\mathbf{A}_I := \{0, 1\}^I$.

Boolean algebra \cong algebra of **continuous** functions from a Stone space X to $\{0, 1\}$ with pointwise operations.

2. Priestley duality (1969): $\mathbf{L} = \{0, 1\}$.

Space: a Stone space.

Constraints: the order, i.e. $\mathbf{A}_I := \{\text{order-preserving } I \rightarrow \{0, 1\}\}$.

Bounded distributive lattice \cong algebra of **order-preserving continuous** functions from a Priestley space X to $\{0 < 1\}$.

" $I \subseteq_2 X$ " means " I is a subset of X of cardinality at most 2".

Definition

A Priestley \mathbf{L} -space consists of a Stone space X and, for each $I \subseteq_2 X$, of a subalgebra \mathbf{A}_I of \mathbf{L}^I s.t.

1. (**Separation**) For all distinct $x, y \in X$, there is $f \in \mathbf{A}_{\{x,y\}}$ s.t. $f(x) \neq f(y)$.
2. (**Local-to-global extension**) For all $I \subseteq_2 X$, every $f: I \rightarrow \mathbf{L}$ in \mathbf{A}_I has a continuous extension $g: X \rightarrow \mathbf{L}$ (with \mathbf{L} discrete) that satisfies all constraints, i.e., s.t., for all $J \subseteq_2 X$, $g|_J \in \mathbf{A}_J$.

Separation



Anti-symmetry of the order.

Local-to-global extension



for $x \not\leq y$ there is a clopen upset U s.t. $x \in U, y \notin U$.

The algebra associated to a Priestley \mathbf{L} -space is

$$\{f: X \rightarrow \mathbf{L} \mid f \text{ is continuous and satisfies all constraints}\}.$$

\mathbf{K} := Kleene algebra $\{0, U, 1\}$.

Example of a \mathbf{K} -Priestley space:

- ▶ $\alpha\mathbb{N}$:= one-point compactification of the discrete space \mathbb{N} .
- ▶ Set $\mathbf{A}_\emptyset := \mathbf{K}^\emptyset$.
- ▶ For every $x \in \alpha\mathbb{N}$ set $\mathbf{A}_{\{x\}} := \mathbf{K}^{\{x\}}$.
- ▶ For distinct $x, y \in \mathbb{N}$, set $\mathbf{A}_{\{x,y\}} := \mathbf{K}^{\{x,y\}}$.
- ▶ For $x \in \mathbb{N}$ set

$$\mathbf{A}_{\{x,\infty\}} := \{f: \{x, \infty\} \rightarrow \mathbf{K} \mid f(\infty) = U \text{ or } f \equiv 0 \text{ or } f \equiv 1\}.$$

Dual Kleene algebra: algebra of functions $f: \alpha\mathbb{N} \rightarrow \mathbf{K}$ such that $f \equiv 0$ or $f \equiv 1$ or $(f(\infty) = U$ and $f(x) = U$ for all but finitely many $x \in \mathbb{N}$).

Duality theorem

Hypotheses on the generating algebra \mathbf{L} :

0. \mathbf{L} is finite.
1. \mathbf{L} has a lattice reduct.
2. \mathbf{L} has two distinct constants and every subalgebra of \mathbf{L} is simple.
3. For each subalgebra \mathbf{A} of \mathbf{L} , the inclusion $\mathbf{A} \hookrightarrow \mathbf{L}$ is the unique homomorphism from \mathbf{A} to \mathbf{L} .

Theorem

Suppose \mathbf{L} satisfies (0–3). The following categories are dually equivalent.

- ▶ $\mathbf{ISP}(\mathbf{L})$ (with homomorphisms as morphisms)
- ▶ the category of Priestley \mathbf{L} -spaces (and appropriate morphisms).

Our generalisation

Our main result: we give an analagous duality, where the generating algebra \mathbf{L} is allowed to be infinite.

Our motivation:

- ▶ representation of some *MV-algebras*, i.e. the algebras of Łukasiewicz many-valued logic: $\{0, 1\}$ is replaced by $[0, 1]$.
- ▶ representation of some *positive MV-algebras*, i.e. the negation-free version of MV-algebras.

Hypotheses on \mathbf{L} :

- ~~0. \mathbf{L} is finite.~~
1. \mathbf{L} has a lattice reduct, i.e. \mathbf{L} has two terms \vee and \wedge that satisfy the axioms of lattices.
2. \mathbf{L} has two distinct constants and every subalgebra of \mathbf{L} is simple (i.e. the only congruences of a subalgebra \mathbf{A} are the diagonal and $\mathbf{A} \times \mathbf{A}$).
3. For each subalgebra \mathbf{A} of \mathbf{L} , the inclusion $\mathbf{A} \hookrightarrow \mathbf{L}$ is the unique homomorphism from \mathbf{A} to \mathbf{L} .

Theorem (Main result)

Suppose \mathbf{L} satisfies (1–3). The following categories are dually equivalent:

- ▶ the category of algebras $\mathbf{A} \in \mathbf{ISP}(\mathbf{L})$ s.t. for each $a \in \mathbf{A}$ the set

$$\{h(a): h: \mathbf{A} \rightarrow \mathbf{L} \text{ homomorphism}\}$$

is finite.

- ▶ the category of Priestley \mathbf{L} -spaces.

An equivalent characterization for the algebras: there is a set X and an embedding of \mathbf{A} into the algebra of functions $X \rightarrow \mathbf{L}$ with finite image.

E.g.: any finite product of \mathbf{L} , and any subalgebra of it.

Thank you!