Natural dualities with an infinite dualizing object

Marco Abbadini. University of Birmingham, UK.

Work in progress with Adam Přenosil.

AAA104, Blagoevgrad, Bulgaria, 10 February 2024.

1936: Stone's representation theorem for Boolean algebras:

Every Boolean algebra $\langle B; \lor, \land, 0, 1, \neg \rangle$ is isomorphic to a field $\langle \mathcal{A}; \cup, \cap, \varnothing, X, X \setminus (-) \rangle$ of subsets of some set X.

Stone gave a uniform way to represent each Boolean algebra as a field of sets: for a Boolean algebra **B**, Stone constructed a topological space X_{B} , and

$$\begin{split} \mathbf{B} &\cong \{ \text{clopen } (= \text{closed} + \text{open}) \text{ subsets of } X_{\mathbf{B}} \} \\ &\cong \{ f \colon X_{\mathbf{B}} \to \{0,1\} \mid f \text{ is continuous} \}. \end{split}$$

This gives a 1:1 correspondence between Boolean algebras and Stone spaces (a.k.a. profinite spaces or Boolean spaces).

Stone space := compact Hausdorff spaces where clopens separate points.

E.g.: finite discrete spaces, one-point compactifications of discrete spaces.

In fact, Stone proved much more than a representation theorem:

 $\mathsf{Bool}\cong\mathsf{Stone}^{\operatorname{op}}.$

Homomorphisms $A \to B$ of Boolean algebras correspond to continuous functions $X_B \to X_A$.

Algebraic questions in Bool can be answered by translating them into (often simpler) questions in Stone.

- 1. While a coproduct $\mathbf{A} + \mathbf{B}$ of Boolean algebras is often difficult to describe, its dual, $X_{\mathbf{A}+\mathbf{B}}$, is simply the cartesian product $X_{\mathbf{A}} \times X_{\mathbf{B}}$.
- 2. The free Boolean algebra (= Lindenbaum-Tarski algebra) on κ generators has a simple dual description: 2^{κ} .
- 3. Congruences on a Boolean algebra dually correspond to closed subspaces.
- 4. The duality is "logarithmic": a product A × B of Boolean algebras corresponds to a sum of the corresponding spaces: X_{A×B} ≅ X_A + X_B. For example, the dual of a Boolean algebra of 64 = 2 × 2 × 2 × 2 × 2 × 2 elements is the discrete space of 6 = 1 + 1 + 1 + 1 + 1 + 1 elements.

[1969, Priestley]: the category of bounded distributive lattices $\langle B; \lor, \land, 0, 1 \rangle$ is dually equivalent to the category of Priestley spaces.

 $\mathsf{BDL} \cong \mathsf{Priestley}^{\mathrm{op}}$.

Priestley space := **Stone space** equipped with a **partial order** such that if $x \not\leq y$ then there is a clopen upset *U* such that $x \in U$ and $y \notin U$.

Every bounded distributive lattice is isomorphic to the lattice of clopen upsets of a Priestley space X. Equivalently, to the lattice of order-preserving continuous functions from X to $\{0, 1\}$. Dualities for various other classes of algebras (usually stemming from nonclassical logic) were found: distributive lattices, semilattices with 1 (Hofmann-Mislove-Stralka duality), relative Stone algebras, de Morgan algebras, median algebras, Stone algebras, double Stone algebras, Kleene algebras.

Benefits similar to those for Stone duality apply.

These dualities follow a similar pattern: there is a "special" algebra ${\sf L}$ (for logics, ${\sf L}$ is the algebra of truth values) and a duality is given for

$\mathbb{ISP}(L).$

Examples:

- Boolean algebras = $\mathbb{ISP}(\{0,1\})$ (by the Ultrafilter Lemma).
- Bounded distributive lattices = $\mathbb{ISP}(\{0,1\})$.

Davey & Werner (1983) showed that many of these follow from a general result: duality for $\mathbb{ISP}(L)$ where L is a finite algebra with a near unanimity term. (E.g.: L has a lattice reduct.)

Example: for $L = \{0, 1\}$ we get Stone/Priestley dualities.

These dualities are called "natural dualities".

We present the restriction of this theory under some stronger hypotheses on L:

- 0. L is finite.
- 1. L has a majority term (e.g., L has a lattice reduct).
- 2. L is finitely subdirectly irreducible (e.g. L has two distinct constants and every subalgebra of L is simple).
- 3. For each subalgebra A of L, the inclusion $A \hookrightarrow L$ is the unique homomorphism from A to L.

For example, the Boolean algebra $\mathbf{L} = \{0, 1\}$.

An algebra in $\mathbb{ISP}(L)$ is represented via a

- 1. Stone space X, together with
- 2. a family of "at-most-binary constraints", i.e., for each subset $I \subseteq X$ of cardinality ≤ 2 , a subalgebra A_I of L^I .

satisfying certain compatibility conditions.

Space X +constraints $(\mathbf{A}_I)_I$ \rightsquigarrow Algebra of continuous functions $f: X \to \mathbf{L}$ (with \mathbf{L} topologically discrete) satisfying all constraints, i.e. such that, for all $I \subseteq X$ with $|I| \leq 2$, $f|_I \in \mathbf{A}_I$. 1. Stone duality for Boolean algebras (1936): $L = \{0, 1\}$. Space: a Stone space.

Constraints: "none", i.e. $\mathbf{A}_I \coloneqq \{0, 1\}^I$.

Boolean algebra \cong algebra of continuous functions from a Stone space X to $\{0,1\}$ with pointwise operations.

2. Priestley duality (1969): $\mathbf{L} = \{0, 1\}$.

Space: a Stone space.

Constraints: the order, i.e. $\mathbf{A}_I := \{ \text{order-preserving } I \rightarrow \{0, 1\} \}.$

Bounded distributive lattice \cong algebra of order-preserving continuous functions from a Priestley space X to $\{0 < 1\}$.

" $I \subseteq_2 X$ " means "I is a subset of X of cardinality at most 2".

Definition

A *Priestley* L-space consists of a Stone space X and, for each $I \subseteq_2 X$, of a subalgebra A_I of L^I s.t.

- 1. (Separation) For all distinct $x, y \in X$, there is $f \in \mathbf{A}_{\{x,y\}}$ s.t. $f(x) \neq f(y)$.
- 2. (Local-to-global extension) For all $I \subseteq_2 X$, every $f: I \to L$ in A_I has a continuous extension $g: X \to L$ (with L discrete) that satisfies all constraints, i.e., s.t., for all $J \subseteq_2 X$, $g|_J \in A_J$.

Separation \longleftrightarrow Anti-symmetry of the order.Local-to-global extension \longleftrightarrow for $x \nleq y$ there is a clopen
upset U s.t. $x \in U, y \notin U$.

The algebra associated to a Priestley $\ensuremath{\textbf{L}}\xspace$ is

 $\{f: X \to \mathbf{L} \mid f \text{ is continuous and satisfies all constraints}\}.$

Hypotheses on the generating algebra $\ensuremath{\mathsf{L}}$:

- 0. L is finite.
- 1. L has a majority term.
- 2. L is finitely subdirectly irreducible.
- 3. For each subalgebra A of L, the inclusion $A \hookrightarrow L$ is the unique homomorphism from A to L.

Theorem

Suppose L satisifes (0–3). The following categories are dually equivalent.

- ▶ ISP(L) (with homomorphisms as morphisms)
- the category of Priestley L-spaces (and appropriate morphisms).

Our generalisation

Our main result: we give an analagous duality, where the generating algebra ${\bm L}$ is allowed to be infinite.

Our motivation:

- representation of some *MV-algebras*, i.e. the algebras of Łukasiewicz many-valued logic: {0,1} is replaced by [0,1].
- representation of some *positive MV-algebras*, i.e. the negation-free version of MV-algebras.

Hypotheses on L:

- 0. L is finite.
- 1. L has a majority term.
- 2. L is finitely subdirectly irreducible.
- 3. For each subalgebra A of L, the inclusion $A \hookrightarrow L$ is the unique homomorphism from A to L.

Theorem (Main result)

Suppose L satisifes (1–3). The following categories are dually equivalent:

▶ the category of algebras $A \in ISP(L)$ s.t. for each $a \in A$ the set

 ${h(a): h: \mathbf{A} \rightarrow \mathbf{L} \text{ homomorphism}}$

is finite.

the category of Priestley L-spaces.

An equivalent characterization for the algebras: there is a set X and an embedding of **A** into the algebra of functions $X \rightarrow L$ with finite image. E.g.: any finite product of **L**, and any subalgebra of a finite product of **L**.

For $\boldsymbol{\mathsf{L}}$ finite, this coincides with the theorem mentioned in the finite case.

Thank you!