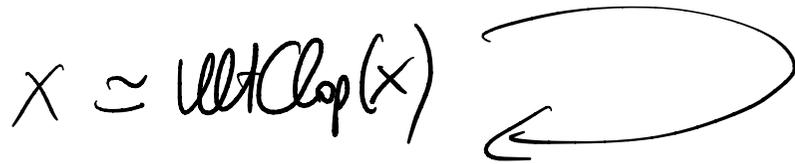
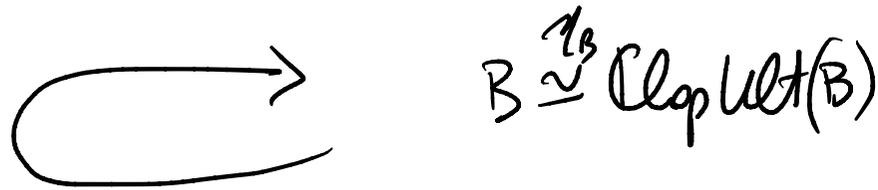
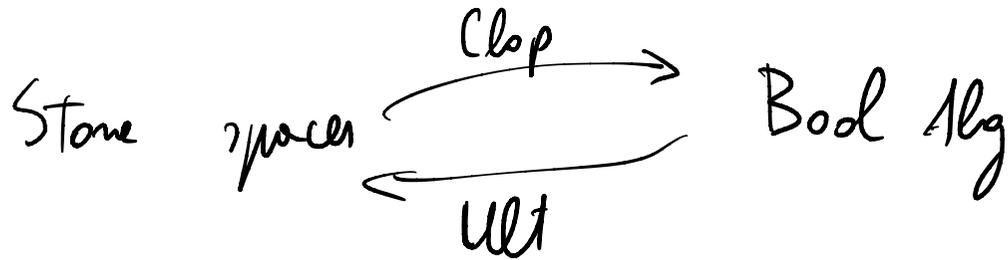
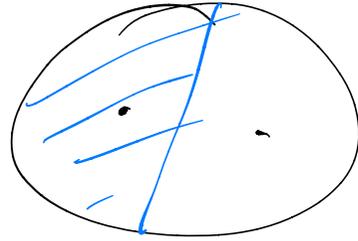


LAST TIME:

2026-02-25

4th lecture

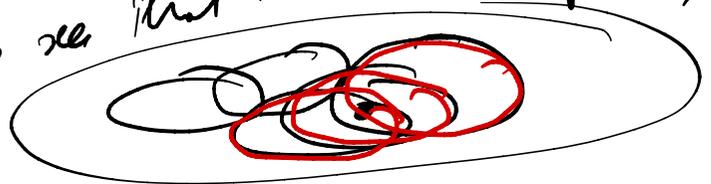
Stone spaces: $\left\{ \begin{array}{l} \rightarrow \text{totally separated (= clopens separate the points)} \\ \rightarrow \text{compact} \end{array} \right.$



THM Let X be a Stone space. The function

$$\begin{aligned} \varepsilon_x: X &\longrightarrow \text{Ult}(\text{Clop}(X)) \\ x &\longmapsto \{A \in \text{Clop}(X) \mid x \in A\} \end{aligned}$$

Easy to see that is an ultrafilter \mathcal{X}



is a homeomorphism

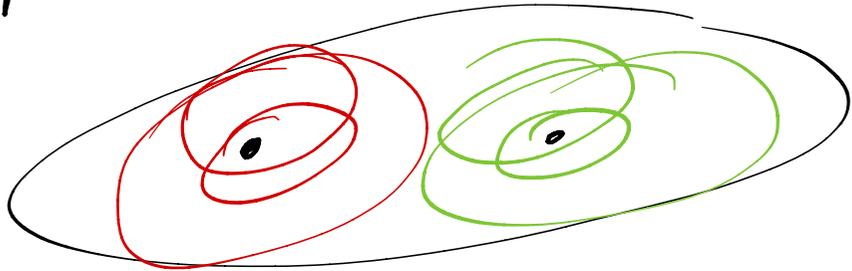
PROOF

Enough \rightarrow CONTINUOUS (IMMEDIATE)
 \rightarrow INJ. (TOTAL SEPARATION)
 \rightarrow SURJ. (COMPACTNESS)

CONTINUITY: enough to prove that the preimage of basic open set

is open: $E_x^{-1}[\underbrace{\mathcal{N}_B(A)}_{\uparrow \text{clap}(x)}] = A$ is open.
 basic open set of $\text{Ult}(x)$

INJ: Let $x \neq y$. By total separation, there is a clopen A s.t. $x \in A$ and $y \notin A$.

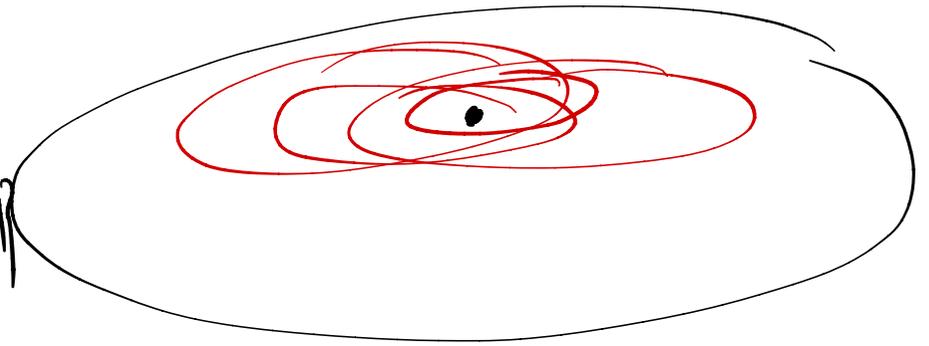


$A \in E_x(x)$
 $\notin E_x(y)$

SURT: Let $\mathcal{U} \in \mathcal{U}lt(\text{Clopt}(X))$

GOAL: find $x \in X$ s.t.

$$\mathcal{U} = \varepsilon_x(x) = \{A \in \text{Clopt}(X) \mid x \in A\}$$



CLAIM: $\bigcap_{A \in \mathcal{U}} A \neq \emptyset$

The family $\{A\}_{A \in \mathcal{U}}$ has the FIP.
This proves the claim.

$$\Rightarrow \exists x \in \bigcap_{A \in \mathcal{U}} A$$

Then it is easily seen that

$$\mathcal{U} = \varepsilon_x(x), \text{ i.e.}$$

$$\mathcal{U} = \{A \in \text{Clopt}(X) \mid x \in A\}$$

EASY

\Rightarrow We prove the contrapositive:

$$A \notin \mathcal{U} \Rightarrow x \notin A \Rightarrow x \in X \setminus A \Rightarrow x \notin A$$

REF

Let X be a compact space.

Let $\{A_i\}_{i \in I}$ be a family of closed sets.

If $\{A_i\}_{i \in I}$ has the FINITE INTERSECTION PROPERTY (FIP): i.e. for every finite $J \subseteq I$, $\bigcap_{i \in J} A_i \neq \emptyset$.

then $\bigcap_{i \in I} A_i \neq \emptyset$

If

- each $A_i \neq \emptyset$
- $\{A_i\}_{i \in I}$ is closed under finite intersection

then it has the FIP.

Bool. Alg \iff Stone spaces

From $B = \text{Clop Ult}(B)$ we deduce

If B is finite, then B is isomorphic to a power set of a finite set.

Indeed: B fin \Rightarrow $\text{Ult}(B)$ is finite

\Downarrow

$\text{Ult}(B)$ is a finite discrete space

\Downarrow

$$\text{Clop}(\text{Ult}(B)) = \mathcal{P}(\text{Ult}(B))$$

\Downarrow
 B

If two ^{FINITE} Bool. alg. have the same cardinality, they are isom.
 $\mathcal{P}(X) \quad \mathcal{P}(Y) \quad 2^n = 2^m \Rightarrow n = m$

Let $L = \{p_0, p_1, p_2, \dots\}$. We will prove

"There is no propositional theory that encodes the information 'exactly one among all p_i 's holds"

$$\forall i \neq j \neg(p_i \wedge p_j)$$

$$\bigvee_{i \in \mathbb{N}} p_i$$

REWRITE MORE FORMALLY:

'There is no propositional theory T such that

$$\text{Mod}(T) = \{v: L \rightarrow \{0,1\} \mid \text{exactly one among all } v(p_i), i \in \mathbb{N}\}$$

REMARK

For any prop. theory T in \mathcal{L} , $\text{Mod}(T)$ is a Stone space,
with clopens: for every formula $\varphi \in \text{Form}(\mathcal{L})$

$$\{\underbrace{v \in \text{Mod}(T)}_{\bar{v}(\varphi)=1} \mid v \models \varphi\}$$

SKETCH of PROOF

$$\text{Mod}(T) \longleftrightarrow \text{Ult}\left(\text{Form}(\mathcal{L}) / \equiv_T\right)$$

$$v \longmapsto \{[\varphi] \mid v \models \varphi\}$$

$$\begin{array}{l} \mathcal{L} \rightarrow \{0, 1\} \\ \mu \mapsto \begin{cases} 1 & [\varphi] \in U \\ 0 & [\varphi] \notin U \end{cases} \end{array}$$

If BwOC there was a theory T with the above property

$$\text{Mod}(T) = \bigcup_{m \in \mathbb{N}} \underbrace{\{v \in \text{Mod}(T) \mid v \models \rho_m\}}_{\text{clopens}}$$

By compactness, there would be $I \subseteq \mathbb{N}$ finite s.t.

$$\text{Mod}(T) = \bigcup_{m \in I} \{v \in \text{Mod}(T) \mid v \models \rho_m\}$$

Take $m \in \mathbb{N} \setminus I$ $v_m: \mathcal{L} \rightarrow 2$
 $i \mapsto \begin{cases} 1 & i = m \\ 0 & \text{if not} \end{cases}$

$m \in I: v_m \not\models \rho_m$

One can prove, as a consequence of the fact that the spaces of models are compact.

[THM (Compactness theorem of classical prop. logic)

Let T be a prop. theory in a prop. language L , $\varphi \in \text{Form}(L)$

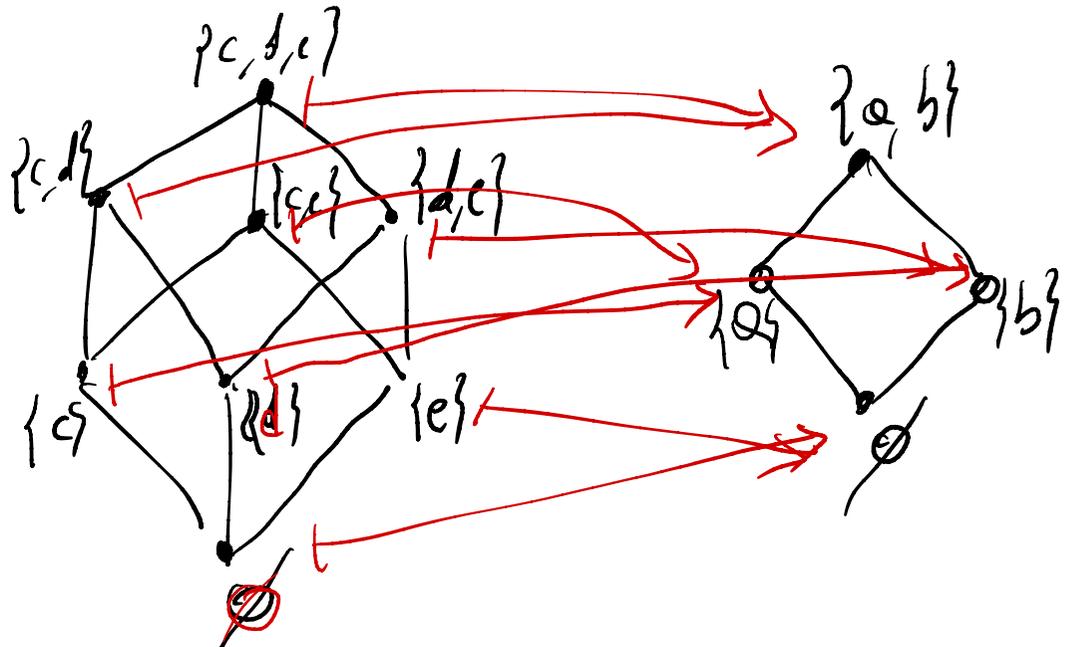
Suppose $T \models \varphi$ (i.e. for every model ν of T , $\nu \models \varphi$
 $(\nu(\varphi) = 1)$)

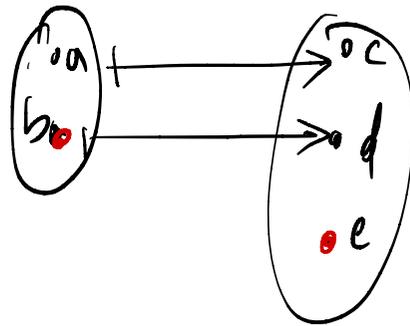
Then, there is a finite $T' \subseteq T$ s.t. $T' \models \varphi$

We skip the proof.

(Logical meaning of Bool. hom.:
 interpretation of formulas)

Stone proved a repr. of Bool. hom.

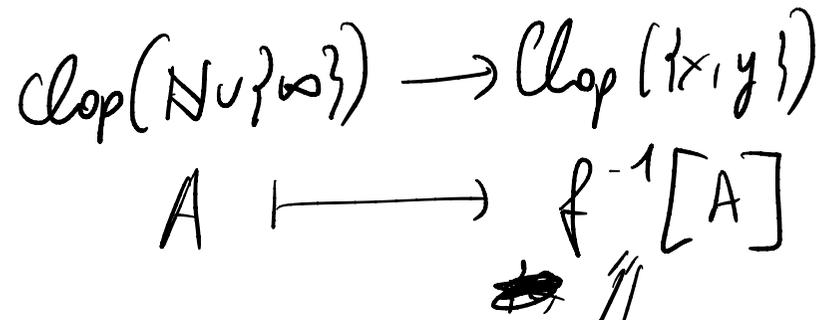
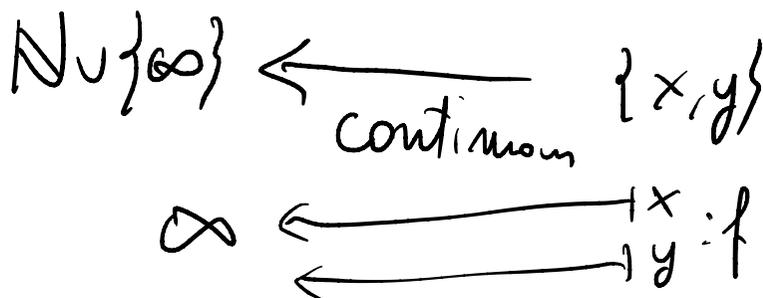
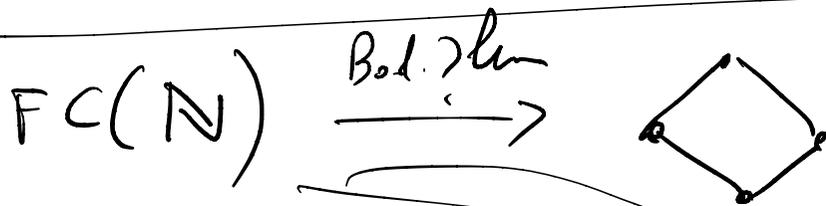


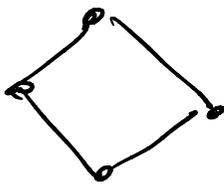


$$f: \{a, b\} \rightarrow \{c, d, e\}$$

$$P(\{c, d, e\}) \rightarrow P(\{a, b\})$$

$$z \mapsto f^{-1}[z]$$



$FC(\mathbb{N}) \rightarrow$ 

$\{x, y\}$ $\in A$
 \emptyset $\notin A$

$F \mapsto \begin{cases} 1 \\ 0 \end{cases}$ F cof.
 F finite

$BA :=$ category of Bod. algs and Bod. hom.

$Stone :=$ " " Stone spaces and cont. maps

THM (Stone duality for Bod. algs.)

The categories BA and $Stone$ are dually equivalent, i.e.

$$BA \simeq Stone^{op}$$

Roughly speaking \rightarrow "bijection on objects"

\rightarrow "bijection on morphisms, but reversing the direction!"

Function:

$$\text{Clop}: \text{Stone}^{\text{op}} \longrightarrow \text{BA}$$

$$X \longmapsto \text{Clop}(X)$$

$$\begin{array}{c} \text{Stone} \\ \uparrow f \\ X \end{array}$$

$$\text{Clop}(Y)$$

$$\downarrow \\ \text{Clop}(X)$$

$$f^{-1}[-]: Z \longmapsto f^{-1}[z]$$

$$\text{Uelt}: \text{BA} \longrightarrow \text{Stone}^{\text{op}}$$

$$B \longmapsto \text{Uelt}(B)$$

A

$f \downarrow$

B

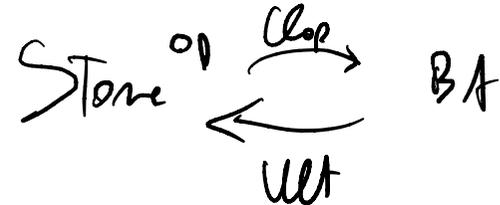
\cong

$\text{Uelt}(A)$ $f^{-1}[u]$

$\uparrow f^{-1}[e]$

$\text{Uelt}(B) \ni u$

NAT. ISO :



$$\text{UNIT: } \eta_B : B \rightarrow \text{Clop Ult}(B)$$

$$f \downarrow \simeq \downarrow \text{Clop Ult}(f)$$

$$A \rightarrow \text{Clop Ult}(B)$$

it is an ISO, as we proved ~~some~~ some times.

$$\text{COUNIT: } \epsilon_X : X \rightarrow \text{Ult Clop}(X)$$

$$f \downarrow \simeq \downarrow \text{Ult Clop}(f)$$

$$Y \rightarrow \text{Ult Clop}(X)$$

NAT. TRANSF.

it is an iso because of the first theorem of today.



It can be easily checked that it is an adj. (e.g., with the triangle identities) \square

GUIDING QUESTION:

given A, B Bool. algs, in the coproduct map

$a \neq 1 \mapsto \neq 1$
 $A \rightarrow A + B$
 injective?

Every equat. defined class of algebras
 in \rightarrow complete (EASY: limits are in Set)
 \hookrightarrow cocomplete (HARD: colimits NOT as in Set)

$2 + 2 = 2$

product map
 $X \times Y \rightarrow X$
 surj.

B/A	Stars
2	{*}
+	x
2	{*}
"	"
2	{*}

1. Explain the logical reading of the quest.
 2. ~~Ex~~ Limits in Stone are computed in $\hat{\mathcal{L}}$
 3. $\text{inj} \leftrightarrow \text{surj.} \quad \text{surj} \leftrightarrow \text{inj.}$
 4. Solve The problem in Stone spaces, get back to BA (and log_2)
-

$$1. \quad A \simeq \text{Form}(L_1) \Big/ \equiv_{T_1} \quad B \simeq \text{Form}(L_2) \Big/ \equiv_{T_2}$$

Let T_1, T_2 be prop. theories in language L_1, L_2 , resp.

For all $\varphi \in \text{Form}(L_1)$ s.t. $T_1 \not\models \varphi$ (i.e. there is a model \mathcal{V} of T_1)
 s.t. $\mathcal{V}(\varphi) = 0$

$$\underbrace{T_1 \cup T_2}_{\text{in } L_1 \cup L_2} \not\models \varphi \quad ?$$

$$\mathcal{L}_1 = \{p, q\}$$

$$\tau_1 = \{p, v, q\}$$

$$\tau_2 \neq \{p, v, q\}, \text{ because } v: \begin{array}{l} \mathcal{L} \mapsto 2 \\ p \mapsto 1 \\ q \mapsto 0 \end{array}$$

$$\mathcal{L}_2 = \{r\}$$

$$\tau_2 = \{r, v\}$$

$$\tau_1 \perp \tau_2 \neq \underline{\{p, v, q\}}, \text{ because } v: \begin{array}{l} p \mapsto 0 \\ q \mapsto 0 \\ r \mapsto 0 \end{array}$$

$\{p, v, q, r\}$

$$\text{Fon}(\mathcal{L}_1) \Big/ \equiv_{\tau_1} + \text{Fon}(\mathcal{L}_2) \Big/ \equiv_{\tau_2} = \text{Fon}(\mathcal{L}_1 \perp \mathcal{L}_2) \Big/ \equiv_{(\tau_1 \perp \tau_2)}$$

$$\text{Mod}(\tau_1 \perp \tau_2) \Leftrightarrow \text{Mod}(\tau_1) \times \text{Mod}(\tau_2)$$

$$A \rightarrow A \times B$$

$$(X \times Y) \rightarrow X$$

2. LIMITS in Stone: easy: as in Set (as in Top)

Examples: Tern. obj: $\{*\}$

Bin. prod. $X_1 \times X_2$ cart. prod. with the product topology

Arb. prod. $\prod X_i$ cart. prod. with the product topology.

(Tychonoff's theorem: arbitrary product of compact spaces is compact)

\Updownarrow

Ax. of Choice

Equalizers: $Z \hookrightarrow X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

$\text{Eq}(f, g) = \{x \in X \mid f(x) = g(x)\}$, equipped with the subspace topology.

This is a Stone space because:

- Enough to prove that $\{x \in X \mid f(x) = g(x)\}$ is closed in X .

it is the preimage ~~to~~ under

$$X \xrightarrow{\langle f, g \rangle} Y \times Y$$
$$x \longmapsto (f(x), g(x))$$

of the diagonal

$\{(y, y) \mid y \in Y\}$ ~~set~~ which

is closed because

Y is Hausdorff.

(Hausdorff = disjoint closed)

LEMMA:

Every closed subspace of a Stone space is a Stone space (with the subspace topology) \square Ex

THM

The functors

$$\text{Stone} \rightarrow \text{Set}, \quad \text{Stone} \rightarrow \text{Top}$$

preserve limit

BTW,

merely finite coproducts, but not arbitrary coproducts. Also, it does not preserve pushouts and coequalizers.

PROOF See above for the limit

Question from the students:

Does $\text{Stone} \rightarrow \text{Set}$ have a left adjoint?

(It cannot have a right adjoint, because it does not preserve colimits)
YES, it has a left adjoint $\text{Set} \rightarrow \text{Stone}$, called the Stone-Cech compactification: and X is mapped to $\text{ultra}(P(X))$

3.

THM ($\text{inj} \Leftrightarrow \text{surj}, \text{surj} \Leftrightarrow \text{inj}$)

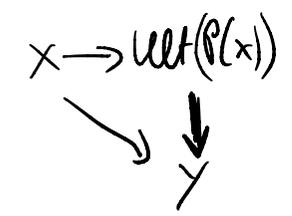
$$A \xrightarrow{\text{inj?}} A \times B$$
$$\boxed{X \times Y} \xrightarrow{\text{surj?}} X$$

Let X, Y be Stone space, and $f: X \rightarrow X$ a cont. map.

- ① f is inj. iff $f^{-1}[-]: \text{Clop}(Y) \rightarrow \text{Clop}(X)$ is surj.
- ② f is surj iff $f^{-1}[-]: \text{Clop}(X) \rightarrow \text{Clop}(X)$ is inj.

If $g: A \rightarrow B$ is a Bod. hom.

- ① g is surj $\Leftrightarrow g^{-1}[-]: \text{ultra}(B) \rightarrow \text{ultra}(A)$ is INJ
- ② g is inj $\Leftrightarrow g^{-1}[-]$ is SURJ.

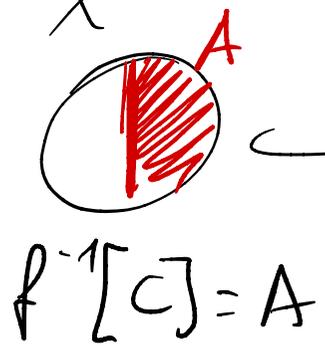


PROOF

① \Rightarrow) f is inj.

Let $A \in \text{Clop}(x)$.

Find a clopen C of Y s.t.



$f^{-1}[C] = A$

$f[A]$

Apply the lemma with

$f[A]$ and $f[X \setminus A]$

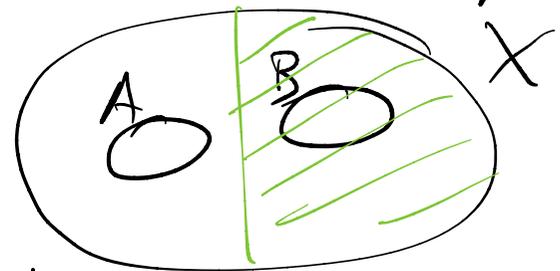
\Rightarrow we have a clopen C s.t.

$f[A] \subseteq C$

$f[X \setminus A] \cap C = \emptyset$

LEMMA

Let X be a Stone space



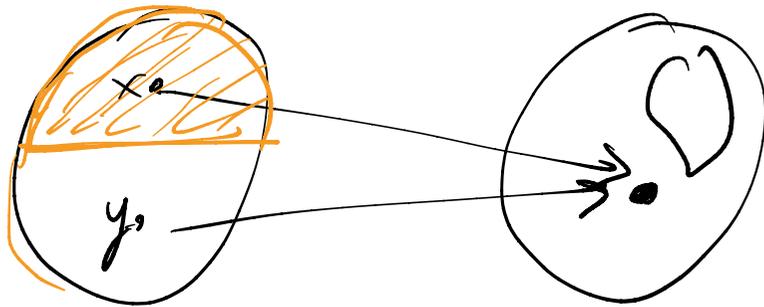
Let A, B closed disjoint subsets,
Then there is a clopen C of X
s.t. $B \subseteq C$ $A \cap C = \emptyset$.

EX

Then $A = f^{-1}[C]$. Then $f^{-1}[-]$ is surj.

\Leftrightarrow) Let us prove the contrapositive

Suppose f not inj'. GOAL: $f^{-1}[-]$ is not surj.



Take any clopen separating x and y . It is not the preimage of any clopen of Y .

② BE CONTINUED...