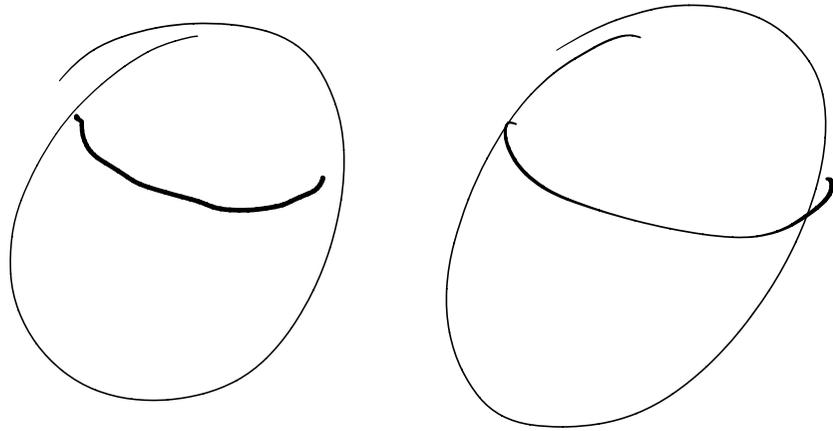


• ERRATA CORRECT: ① Stone \rightarrow Set

preserves finite coproducts

INIT: \emptyset

BIN. COPRODUCT: $X + Y$ is the disjoint union with the disjoint union topology.



② Stone \rightarrow Set

does not preserve coequalizers

(see example ^{few pages} below).

BA is a variety of algebras (= eq. definable).

\Rightarrow complete and cocomplete.

Therefore,
Stone is complete and cocomplete

Stone \rightarrow Set does not preserve arbitrary colimits

$1 + 1 + \dots$

So Stone \rightarrow Set is not a left adjoint.

Does Stone \rightarrow Set have a left adjoint?

$F: \text{Set} \rightarrow \text{Stone}?$

If F exists then it should satisfy

$$F(X) = F\left(\sum_{z \in X} \{*\}\right) = \sum_{z \in X} F(\{*\})$$

\uparrow
Set

$$\text{hom}_{\text{Stone}}(F(\{*\}), Y) \cong \text{hom}_{\text{Set}}(\{*\}, UY) \cong UY$$

\uparrow
Stone

We can take $F(\{*\}) = \{*\}$

Stone \rightarrow Set is nat. isom. to $\text{hom}_{\text{Stone}}(\{*\}, -)$

$$F(X) := \sum_{x \in X} \{*\}$$

$\sum_{n \in \mathbb{N}} \{*\}$ in Stone?

is the dual of

$\prod_{n \in \mathbb{N}} 2 \cong P(\mathbb{N})$ in Bool. alg.

$$\sum_{n \in \mathbb{N}} \{*\} \cong \text{Ult}(P(\mathbb{N}))$$

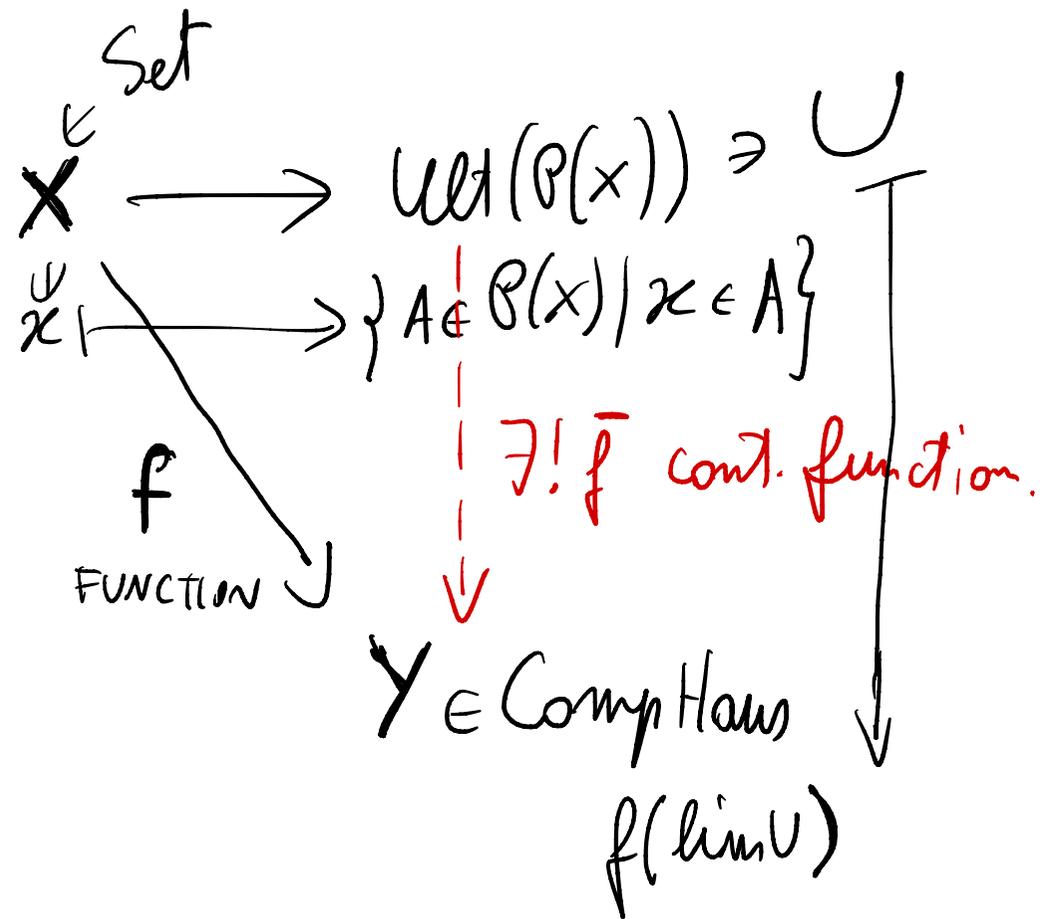
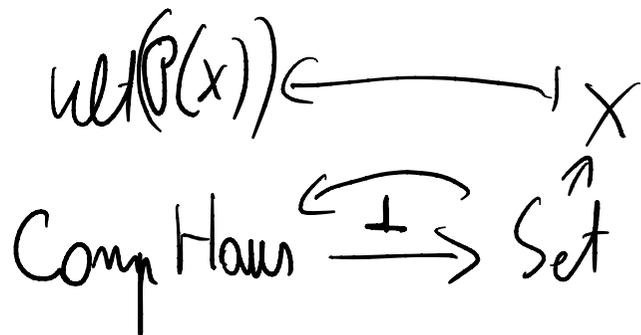
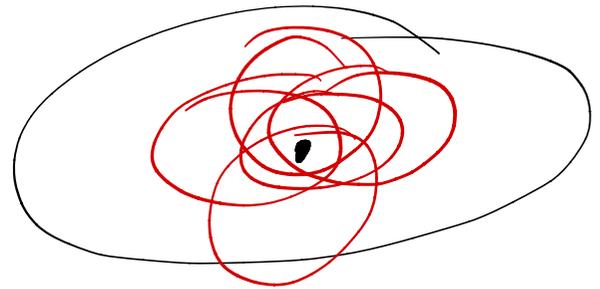
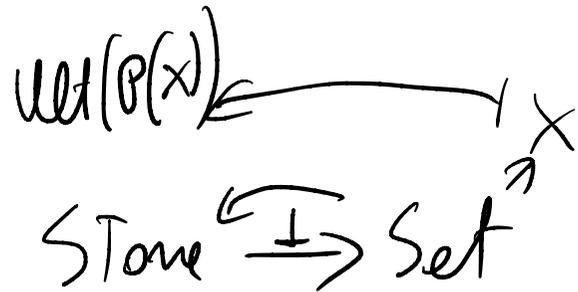
$$\sum_{x \in X} \{*\} \cong \text{Ult}(P(X))$$

(THM: If \mathcal{C} is a cocomplete category, a faithful functor

$$\mathcal{C} \rightarrow \text{Set}$$

has a left adjoint if and only if it is representable

$$\Leftarrow F(x) = \sum_{z \in X} z$$



$\lim U$ is the unique element $z \in X$
s.t. U converges to z .
i.e. every neighbourhood of z belongs to U

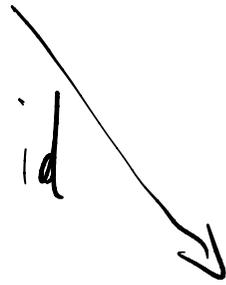


We use it to prove that Stone \rightarrow Set does not preserve
colimits.

CLAIM: for every comp. Hausd. space X there is a Stone space
 Y and a surject. cont. function

$$Y \twoheadrightarrow X$$

$$\overbrace{[0,1]}^{\in \text{Set}} \longrightarrow \text{Ult}(\mathcal{P}([0,1])) \in \text{Stone}$$



$\in \text{Comp Haus}$

$\exists!$ f cont. function

this is surj. because
id is surjective

Let

$$\begin{array}{c} \text{Stone} \\ \downarrow \\ Y \xrightarrow{f} [0,1] \end{array}$$

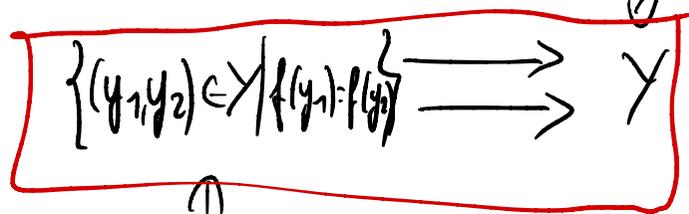
any such map

$Y \times Y$

V_1

Stone

\downarrow



Y

\xrightarrow{f}

$[0,1]$ connected

Stone

one can prove that this is true.

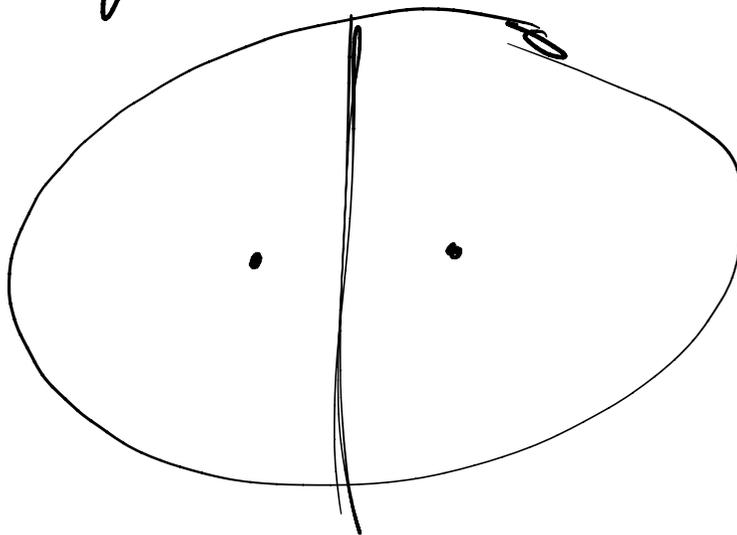


$z \in \text{Stone}$

CLAIM: z is the singleton

The image $g([0,1])$ is connected

The only connected subsets are the singletons



LIMITS AND COLIMITS IN STONE AND BA

Both are complete and cocomplete.

In both cases, limits are computed as in Set.

for BA this is because this holds
for any variety of algebras.
because the forgetful functor

$$BA \longrightarrow \text{Set}$$

is isom.

$$\text{Hom}(\text{Free}_{BA}(I), -)$$

THM ($\text{inj} \Leftrightarrow \text{surj}, \text{surj} \Leftrightarrow \text{inj}$)

Let X, Y be Stone spaces, $f: X \rightarrow Y$ a cont. function

- ① f is $\text{inj} \Leftrightarrow f^{-1}[-]: \text{Clop}(Y) \rightarrow \text{Clop}(X)$ is surjective
- ② f is $\text{surj} \Leftrightarrow f^{-1}[-]: \text{Clop}(Y) \rightarrow \text{Clop}(X)$ is injective

PROOF ① already done.

② \Rightarrow Suppose f surj.

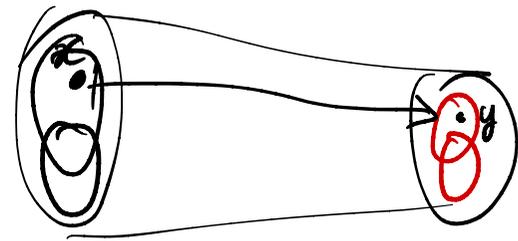
Let $A, B \in \text{Clop}(Y)$ s.t. $A \neq B$

CLAIM: $f^{-1}[A] \neq f^{-1}[B]$.

There is y belonging to A and not B (or vice versa)

By surj., there is x s.t. $f(x) = y$.

$x \in f^{-1}[A] \setminus f^{-1}[B]$, proving the claim.



\Leftarrow) Suppose f is not surj. Let us prove $f^{-1}[-]$ is not inj.

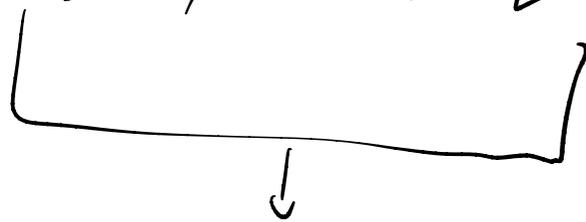
There is $y \in Y \setminus f[X]$

\Rightarrow There is a clopen

C s.t. $y \in C$

$C \cap f[X] = \emptyset$.

$$\Rightarrow f^{-1}[C] = \emptyset \neq f^{-1}[\emptyset]$$



$f^{-1}[-]$ is not inj, because $C \neq \emptyset$.

IDEA:

SURT.
HOM.

$$\text{Form}(\mathcal{L}) / \equiv_T \rightarrow \text{Form}(\mathcal{L}') / \equiv_{T'}$$

$T \subseteq T'$
extend the theory

$$B \hookrightarrow \text{Form}(\mathcal{L}) / \equiv_T$$

"fewer propositions"

INT. CONT. FUNCS

$$\text{Mod}(T') \hookrightarrow \text{Mod}(T)$$

less models.

$$\text{Mod}(T) \twoheadrightarrow Y$$

certain models become equivalent.

$$A \rightarrow A+B \text{ inj?}$$

$$X \times Y \xrightarrow{\pi_X} X \text{ is surj?}$$

$(x,)$ \downarrow
 x

Y is nonempty π_X is surj.
 If $Y = \emptyset$ it is surj iff $X = \emptyset$.

	$Y = \emptyset$	$Y \neq \emptyset$
$X = \emptyset$	SURJ	SURJ
$X \neq \emptyset$	NOT SURJ	SURJ

$$A \rightarrow A+B \text{ inj?}$$

	$B = \{*\}$	$B \neq \{*\}$
$A = \{*\}$	INJ	INJ
$A \neq \{*\}$	NOT INJ	INJ

LOGICAL

$$T_A \quad \varphi \in \mathcal{L}_A$$

$$\text{INT: } \forall \varphi \in \text{Fom}(\mathcal{L}_A) \left(T_A \not\models \varphi \Rightarrow T_A \perp T_B \not\models \varphi \right)$$

Stone

$$\begin{aligned} \text{EPI} &= \text{SURJ} \\ \text{REG. EPI} &= \text{SURJ} \\ \text{MONO} &= \text{INJ} \\ \text{REG. MONO} &= \text{INJ} \end{aligned}$$

Bool. Alg.

$$\begin{aligned} \text{MONO} &\stackrel{?}{=} \text{INJ} \\ \text{REG. MONO} &\stackrel{?}{=} \text{INJ} \\ \text{EPI} &\stackrel{?}{=} \text{SURJ} \\ \text{REG. EPI} &= \text{SURJ} \end{aligned}$$

Both in Stone and BA.

$$\begin{aligned} \text{Epi} &= \text{extra epi} = \text{reg. epi} = \text{SURJ} \\ \text{Mono} &= \text{extra mono} = \text{reg. mono} = \text{INJ} \end{aligned}$$

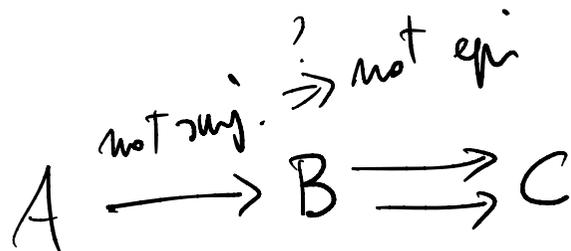
THM. In every variety,
mono = inj.
REG. EPI = SURJ.

EPI of Bool. Alg.

CLAIM: EPI = SURJ.

EPI \Rightarrow SURJ.

SURJ \Rightarrow EPI \checkmark

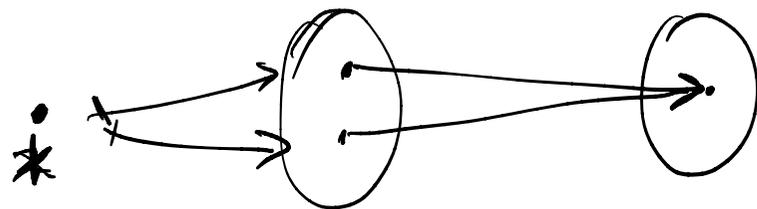


WE USE DUALITY:

We prove in Stone MONO \Rightarrow INJ.

(eq. not inj. implies not mono)

in Stone MONO \Rightarrow INJ.



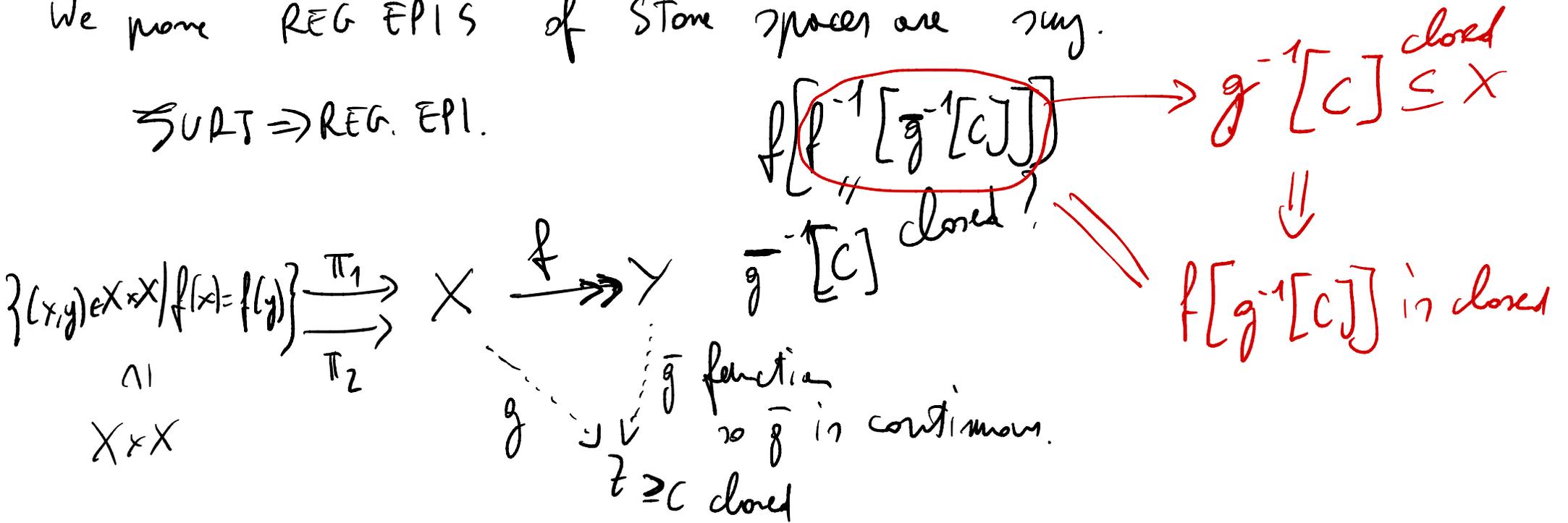
Hence, in BA: EPI \Rightarrow SURJ.

IN LOGICAL TERMS, surjectivity of epimorphisms corresponds to "Beth definability".

REG MONO of Bool. alg.

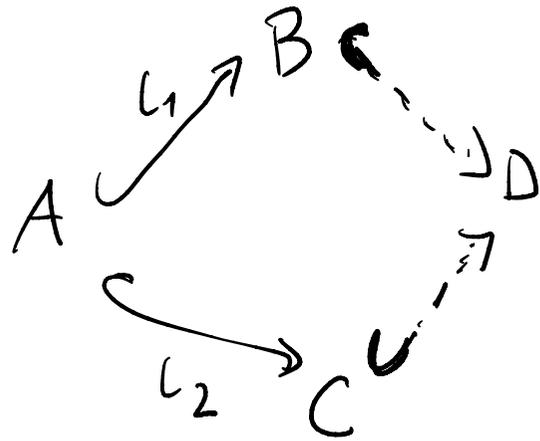
We prove REG EPIS of Stone spaces are surj.

$\text{SURJ} \Rightarrow \text{REG. EPI.}$

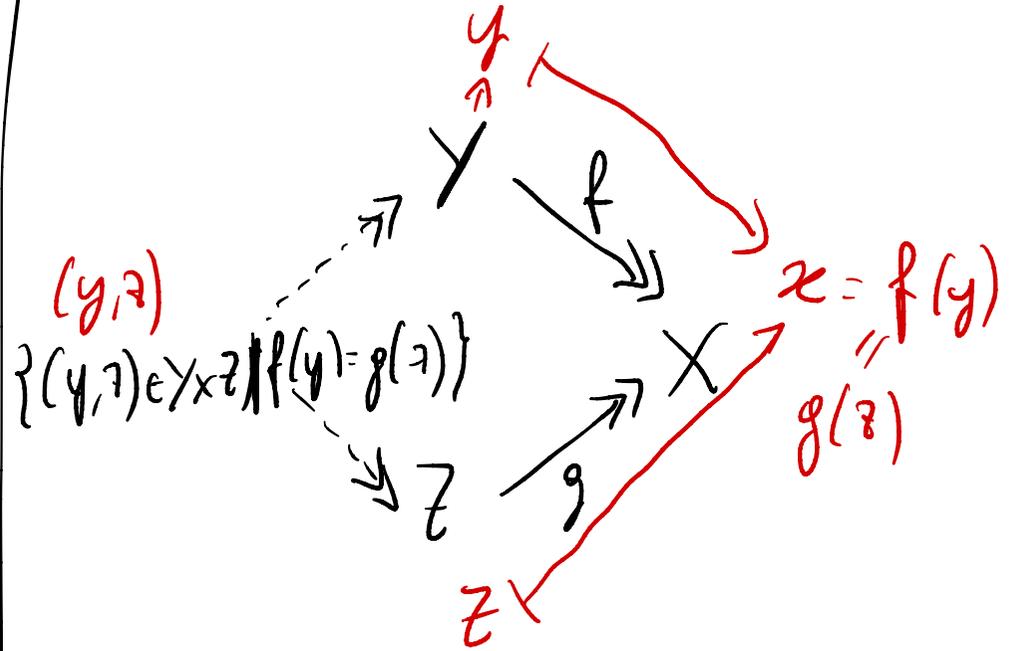


So, f is coreg (π_1, π_2) . So f is reg. epi.

AMALGAMATION (logical Terms: INTERPOLATION)



L_1, L_2 I.N.F. of BA_1



How many formulas, up to logical equivalence, are there with n propositional symbols?

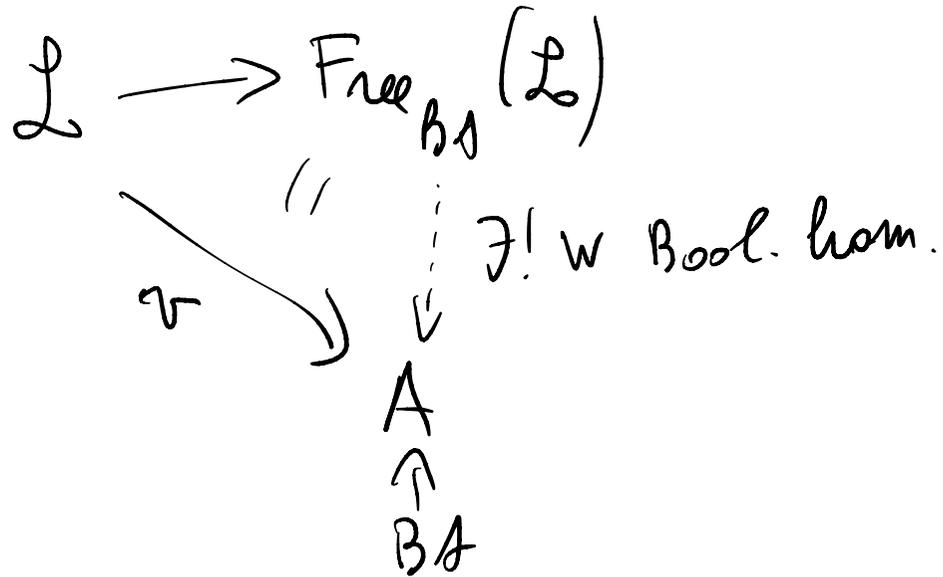
$$n=0 \quad [0], [1]$$

$$n=1$$

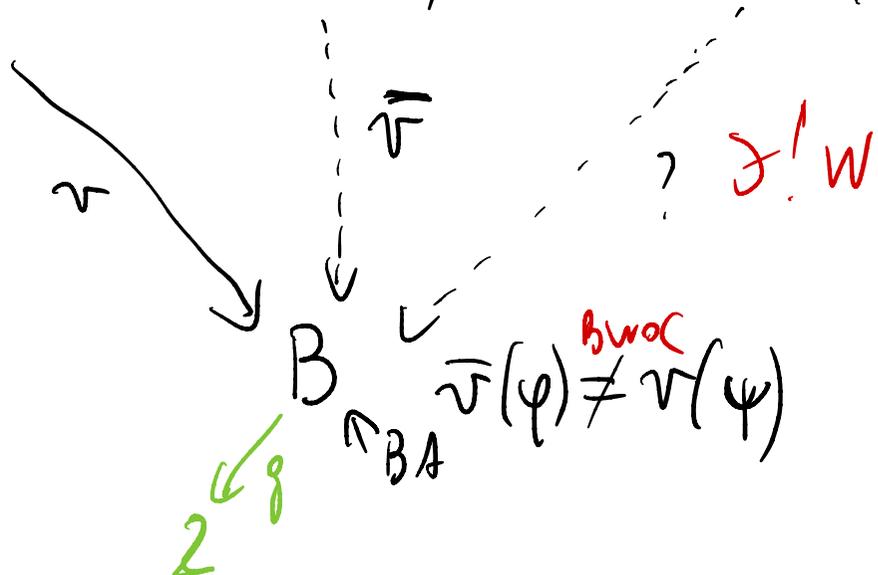
$$[n] \quad [0], [n], [\neg n], [1]$$

$$\text{i.e.: } \left| \text{Form}(\mathcal{L}) / \equiv \right| = ?$$

$$\text{CLAIM: } \text{Form}(\mathcal{L}) / \equiv \simeq \text{Free}_{BA}(\mathcal{L})$$



$$\mathcal{L} \longrightarrow \text{Form}(\mathcal{L}) \xrightarrow{\pi} \text{Form}(\mathcal{L}) / \equiv$$



$\{v, \wedge, \neg, 0, 1\}$

$$\bar{v}(\varphi) \neq v(\varphi) \quad \text{B w o C}$$

Let $\varphi, \psi \in \text{Form}(\mathcal{L})$ s.t. $\pi(\varphi) = \pi(\psi)$, i.e. for every $w: \mathcal{L} \rightarrow 2$
 $\bar{w}(\varphi) = \bar{w}(\psi)$

$$\bar{w}: \text{Form}(\mathcal{L}) \rightarrow 2$$

Does it follow that $\bar{v}(\varphi) = \bar{v}(\psi)$

if not, there is a homo $B \xrightarrow{g} 2$ s.t. $g(\bar{v}(\varphi)) \neq g(\bar{v}(\psi))$

But then, taking $w = v \circ g$, we get a contradiction.

QUESTION:

$$|\text{Free}_{BA}(\mathcal{L})| = ?$$

$$BA \xrightleftharpoons[\cup]{\text{Free}_{BA}} \text{Set}$$

$$\text{Free}_{BA}(\mathcal{L}) = \text{Free}_{BA}\left(\sum_{p \in \mathcal{L}} \{*\}\right) \stackrel{F \text{ is left adj.}}{=} \sum_{p \in \mathcal{L}} \text{Free}_{BA}(\{*\})$$

IT IS DUAL TO THE STONE SPACE

$$\prod_{p \in \mathcal{L}} \text{Ult}\left(\text{Free}_{BA}(\{*\})\right) = \prod_{p \in \mathcal{L}} 2$$

$$\stackrel{|2}{\text{hom}}_{BA}\left(\text{Free}_{BA}(\{*\}), 2\right) \cong \text{hom}_{\text{Set}}(\{*\}, 2) = 2$$

~~IN CO~~ TO SUM UP, the Stone dual of $\text{Free}_{BA}(\mathcal{L})$

is $\prod_{\mu \in \mathcal{L}} 2 = \underline{2^{\mathcal{L}}}$ with the product topology.

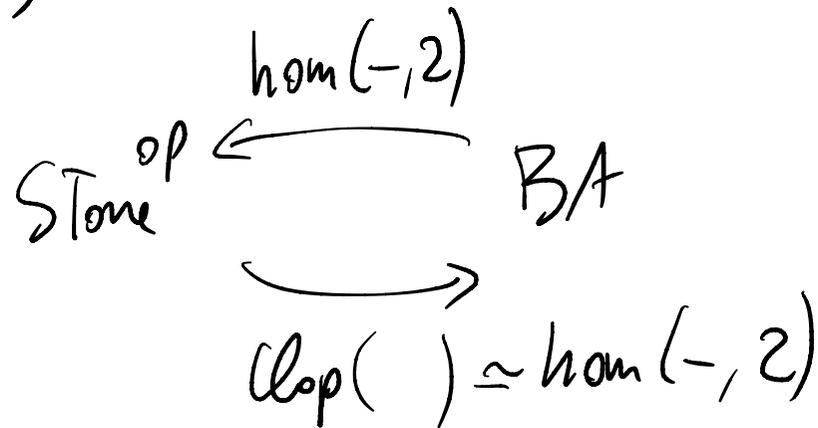
If $|\mathcal{L}| = m$

$$|\underbrace{\text{Free}_{BA}(\mathcal{L})}| = \left| \underbrace{\text{Clop}}_{\text{infinite}} \left(\underbrace{2^{\mathcal{L}}}_{\text{finite}} \right) \right| = \left| \underbrace{\mathcal{P}}_{\text{finite}} \left(\underbrace{2^{\mathcal{L}}}_{\text{finite}} \right) \right| = 2^{2^m}$$

↓ PROPOSITIONAL INTUITIONISTIC LOGIC ~~STON~~

FIRST-ORDER CLASSICAL LOGIC

NATURAL DUALITIES



Set

